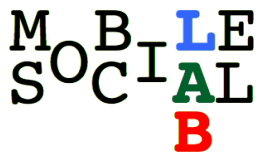


Lecture 12: Infect (2/4)

How do epidemic and gossip reach people?
(i.e., how computer viruses spread?)

COMS 4995-1: Introduction to Social Networks
Wednesday, November 7th



Outline

- * Continuous epidemics, “logistic model”)
- * Discrete epidemics, “graph”
 - Adjacency matrix)
 - SI, SIR model]
 - SIS]
- * Epidemic algorithms

Discrete epidemics

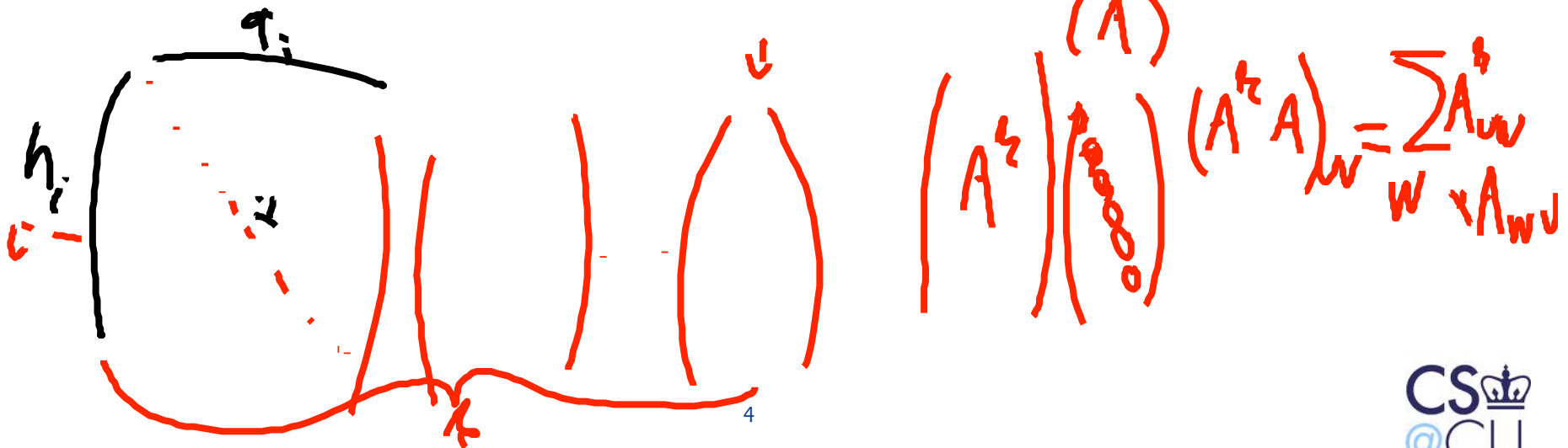
- * Infection only spreads along edges of a given graph
 - to account for connections, closeness among nodes
- * Initial conditions: one or several nodes infected]
- * Challenges :
 - Can a fraction be infected in a large graph?
 - What is the speed of evolution of epidemics?
 - How does it depend on the properties of the graph?]
 - What if some individuals are immune?]

Adjacency matrix

- * A graph $G=(V,E)$ with N nodes
- * Adjacency matrix A is $N \times N$ matrix
 - $A_{uv}=1$ if (u,v) in E , 0 otherwise

$$A_{vu} = 1 \quad (A_{uv} = 0 \quad u=v \quad \text{can be } 1)$$

- * Prop1: for any k , $(A^k)_{uv}$ is $\#\{\text{paths } u \rightarrow v \text{ of length } k\}$



Adjacency matrix

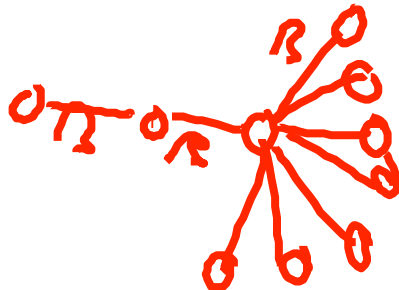
- * Prop2: A is symmetric, with non-negative entries
 - * THM 1 (Perron Frobenius):
 - Eigenvalues of A are all real $\lambda_1 \geq \lambda_2 \geq \dots$
 - λ_1 is the largest in absolute value ($\lambda_1 \geq |\lambda_i|$ for any i)
 - λ_1 has a non-negative eigenvector x_1
- Generally λ_1 is called the **spectral radius** of A , $\rho(A)$

Outline

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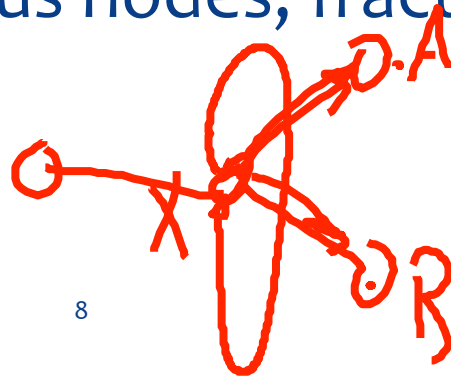
Epidemic Model #1: S→I

- * Model 1: broadcast | *flooding*
 - Node infected at time t infects all its neighbors in $t+1$
 - Within time $D \leq \text{diam}(G)$, all nodes are infected *||||| +*
- * Model 2: gossip
 - Node infects each neighbor with a given rate β
 - Eventually all nodes are infected within $O(D/\beta)$ time)
 - What if rates reduces with degree ? (see later lecture)



Epidemic model #2: $S \rightarrow I \rightarrow R$

- * Model 1: single infection attempts
 - Infected node infects ^{each} neighbors ^{indep} with probability β
 - Many names: “Independent cascade model”, Reed-Frost epidemics, SIR with single slot
- * Model 2: Random infectious period (normalized)
 - Similar (probability to spread is β) but dependencies!
- * Eventually: no infectious nodes, fraction removed



Epidemic model #2: $S \rightarrow I \rightarrow R$

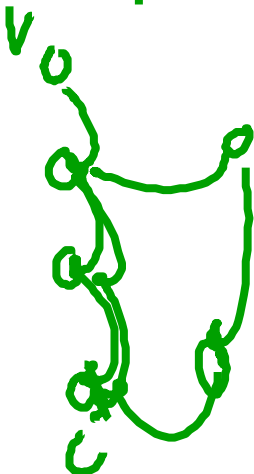
x $Y(\infty)$

- * What is the size of the removed fraction?
- * Thm: Assuming $\beta\rho < 1$, $E[|Y(\infty)|] \leq C \sqrt{N} / (1 - \beta\rho)$
 - $\rho(G)$: largest eigenvalue of G 's adjacency matrix
 - $C = \sqrt{\#\{\text{initial infected population}\}}$
- * If $\beta\rho < 1$ and $C = o(\sqrt{N})$, remove only negligible fraction
- * Proof based on expectation not on independence
 - Similar results hold for model 2

Proof

- * X_v (resp. Y_v) = 1 iff v is infected (resp, inf or recov.)
- * First, we bound $Y_v(t)$ using adjacency matrix

$$\begin{aligned}
 \mathbb{E}[Y(\infty)] &\leq \sum_v \mathbb{P}(Y_v(\infty)=1) \\
 &\leq \sum_{v \text{ infected}} \sum_{t \geq 0} \left[\sum_{u_0=v, u_1, \dots, u_t=v} \beta^t \right] \\
 &\leq \sum_v X_v(0) \sum_{t \geq 0} \beta^t A_{vv}^t \\
 &\leq \sum_v X_v(0) \sum_{t \geq 0} \beta^t A_{vv}^t
 \end{aligned}$$



$y(0) = 0$
 $X(v) = 1$ if v infected, 0 otherwise.
 $X(t) y(t)$

$$|E[y(\infty)]| \leq \sum_{v \in V} \sum_{v \in V} X_v(0) \sum_{t \geq 0} \beta^t A_{vv}^t$$



Proof

- * Second, we use a bounding norm technique



